# CircularTrip and ArcTrip: Effective Grid Access Methods for Continuous Spatial Queries

by

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# THE UNIVERSITY OF NEW SOUTH WALES



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#### Abstract 350 words maximum: (PLEASE TYPE)

A k nearest neighbor guery g retrieves k objects that lie closest to the guery point g among a given set of objects P. With the availability of inexpensive location aware mobile devices, the continuous monitoring of such queries has gained lot of attention and many methods have been proposed for continuously monitoring the kNNs in highly dynamic environment. Multiple continuous queries require real-time results and both the objects and queries issue frequent location updates. Most popular spatial index, R-tree, is not suitable for continuous monitoring of these queries due to its inefficiency in handling frequent updates. Recently, the interest of database community has been shifting towards using grid-based index for continuous queries due to its simplicity and efficient update handling. For kNN queries, the order in which cells of the grid are accessed is very important. In this research, we present two efficient and effective grid access methods, CircularTrip and ArcTrip, that ensure that the number of cells visited for any continuous kNN query is minimum. Our extensive experimental study demonstrates that CircularTrip-based continuous kNN algorithm outperforms existing approaches in terms of both efficiency and space requirement. Moreover, we show that CircularTrip and ArcTrip can be used for many other variants of nearest neighbor queries like constrained nearest neighbor queries, farthest neighbor queries and (k+m)-NN queries. All the algorithms presented for these queries preserve the properties that they visit minimum number of cells for each query and the space requirement is low. Our proposed techniques are flexible and efficient and can be used to answer any query that is hybrid of above mentioned queries. For example, our algorithms can easily be used to efficiently monitor a (k+m) farthest neighbor query in a constrained region with the flexibility that the spatial conditions that constrain the region can be changed by the user at any time.

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### Dedication

I dedicate this thesis to my parents who have always been my nearest and reverse nearest neighbors<sup>1</sup> [KM00] and have been so close to me that I found them with me whenever I needed. It is their unconditional love that motivates me to set higher targets. I also dedicate this thesis to my sisters (Shazia Cheema and Nadia Cheema) and brothers (Muhammad Omer Cheema, Muhammad Umair Cheema and Muhammad Ammar Cheema) who are my nearest surrounders<sup>2</sup> [LLL06] and have provided me with a strong love shield that always surrounds me and never lets any sadness enter inside.

 $<sup>^{1}</sup>$ In a not-so-academic wording, your nearest neighbor is an object that is closest to you and your reverse nearest neighbor is any object for which you are the closest object.

<sup>&</sup>lt;sup>2</sup>Your nearest surrounders are the objects that are nearest to you and surround you completely.

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I feel obliged to say "thanks" to so many people who have been doing so many good things for me but I am afraid the list may go forever. Same as mathematicians write all positive even numbers in very small space as  $\{0, 2, 4, ..., \infty\}$ , below I try to write a shorter acknowledgements. Let k be a positive integer, T be my lifetime, and dist(p,q) be the distance between the hearts of any two people p and q reflecting the love the person p has for q, I feel grateful to all people who have ever been among my k nearest neighbors and/or k reverse nearest neighbors over the time period T.

### Abstract

A k nearest neighbor query q retrieves k objects that lie closest to the query point q among a given set of objects P. With the availability of inexpensive location aware mobile devices, the continuous monitoring of such queries has gained lot of attention and many methods have been proposed for continuously monitoring the kNNs in highly dynamic environment. Multiple continuous queries require real-time results and both the objects and queries issue frequent location updates. Most popular spatial index, R-tree, is not suitable for continuous monitoring of these queries due to its inefficiency in handling frequent updates. Recently, the interest of database community has been shifting towards using grid-based index for continuous queries due to its simplicity and efficient update handling. For kNN queries, the order in which cells of the grid are accessed is very important. In this research, we present two efficient and effective grid access methods, CircularTrip and ArcTrip, that ensure that the number of cells visited for any continuous kNN query is minimum. Our extensive experimental study demonstrates that CircularTrip-based continuous kNN algorithm outperforms existing approaches in terms of both efficiency and space requirement. Moreover, we show that CircularTrip and ArcTrip can be used for many other variants of nearest neighbor queries like constrained nearest neighbor queries, farthest neighbor queries and (k + m)-NN queries. All the algorithms presented for these queries preserve the properties that they visit minimum number of cells for each query and the space requirement is low. Our proposed techniques are flexible and efficient and can be used to answer any query that is hybrid of above mentioned queries. For example, our algorithms can easily be used to efficiently monitor a (k+m) farthest neighbor query in a constrained region with the flexibility that the spatial conditions that constrain the region can be changed by the user at any time.

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